<u>L. LERMAN</u>, S. FERNANDES MEDEIROS, G. BONTEMPI, and O. MARKOWITCH

Université Libre de Bruxelles Faculty of Sciences Department of Computer Sciences Cryptography and Security Service & Machine Learning Group

CARDIS 2013

└─ Cryptography

Context

Cryptography has been used for a long time for confidentiality purposes

Mobile phones

Banks





Side channel attacks

Reduction in cryptography security in real situation Possibility to find the secret key when we focalize on a side channel

- Timing attack (Kocher 1996)
- Electromagnetic attack (Gandolfi, Mourtel & Olivier 2001)
- Power monitoring attack (Kocher, Jaffe & Jun 1999)

Side channel attacks

Power monitoring attack



EM leakage



$$\mathcal{T}^{(Q)} = \left\{ \mathcal{T}^{(Q)}_{(t)} \in \mathbb{R} | t \in [1; n] \right\}$$

 MARTINASEK, Z., ZEMAN, V., TRASY, K.. Simple Electromagnetic Analysis in Cryptography. International Journal of Advances in Telecommunications, Electrotechnics, Signals and Systems, North America, 1, sep. 2012.

Non-profiling attacks

- f is the target function (e.g. SBox) using P and Q
- *L* is the leakage model (e.g. HW)
- D is the distinguisher (e.g. Pearson correlation)

$$\hat{Q} = \arg \max_{Q \in \mathcal{Q}} | D(L(f(P,Q)), T) |$$

└─Side channel attacks

Profiling attacks

$$\hat{Q} = \arg \max_{Q} P(Q|T)$$

$$\hat{Q} = \arg \max_{Q} \frac{P(T|Q) \times P(Q)}{P(T)}$$

$$\hat{Q} = \arg \max_{Q} \hat{P}(T|Q) \times \hat{P}(Q)$$

How to estimate P(T|Q)?

Profiling attacks

Parametric methods

- TA (i.e. $P(T|Q_i) \sim N(\mu_i, \Sigma_i)$) [S. Chari et al. 2002]
- SA (i.e. $P(T|Q_i) \sim N(\mu_i, \Sigma)$) [W. Schindler et al. 2005]
- Non-parametric methods [L. Lerman et al. 2011 & 2013, G. Hospodar et al. 2011, A. Heuser et al. 2012, T. Bartkewitz et al. 2012]
 - SVM
 - RF
 - KNN
- Results in unprotected contexts
 - A ML model is as efficient (and often better) than TA

Countermeasures

- Several countermeasures
 - Masking
 - Hiding
- Several algorithms of masking schemes
 - <u>Boolean</u>, multiplicative, affine masking schemes

└─A machine learning approach against a masked AES

lssues

Are the results of the previous ML works still the same in a protected environment?

How many traces are required

- against a protected device with a ML model compared to a strategy based on TA or SA?
- 2 by a ML model attacking a protected device compared to an unprotected device?
- 2 What is the impact of the number of traces used in the profiling step by a ML model attacking a protected device?

Framework



Lower the error between the correct and the estimated mask values, higher the correlation between the real and the predicted traces for the correct key



- AES-128 protected by the Rotating Sbox (Boolean) Masking scheme (based on table look-up)
- Atmel ATMega-163 smart card
- According to its authors (in a hardware context):
 - Performances and complexity close to unprotected scheme
 - Resistant against several side-channel attacks

A machine learning approach against a masked AES

Models

Profiling attacks

- TA
- SA
- SVM
- RF
- Non-profiling attack
 - CPA on HW(maskedSBox(plaintext \oplus mask \oplus key))



- Public dataset of the DPAContest V4 (updated in October)
- Electromagnetic emission leakages
- First round of AES
- Each trace has 435,002 samples

A machine learning approach against a masked AES

Finding the offset value on traces

 $\rho({}_tT, \text{offset})$ on 1500 traces



Feature selection step: 50 instants highest linearly correlated with the offset value

15 / 20

– Experiments

Model estimation



16 / 20

Model selection results

- Higher the number of traces in the learning set, higher the accuracy
- Higher the number of features, higher the success rates for SVM, RF and SA (except TA)
- The success rates of
 - ML models
 - SVM: 0.88
 - RF: 0.81
 - SA: 0.90
 - **TA: 0.66**

Attacking step



- Unmasked implementation
 - CPA: 16.3 traces in average (5s)
- Masked implementation
 - SVM / CPA: 26 traces in average (20s)
 - SA / CPA: 27.8 traces in average (80s)
 - TA / CPA: 56.4 traces in average (45s)
 - SA: 107 traces in average (180s)

Discussion & Conclusion

Discussion & Conclusion

- (Unprotected) implementation of the Rotating Sbox Masking
 - 26 traces with 20s during the attacking phase
- ML approach outperforms TA in data complexity
- Original SA is less efficient than the new strategy based on SA
- SVM outperforms SA in time complexity
- How to improve the attack ?
 - Increasing the number of points selected in each trace
 - Optimizing the model's parameters

Discussion & Conclusion

Last but not least ...

Official result in the DPAContest V4 :

22 traces with 0.528 seconds

in order to retrieve the secret key of AES-128